- 1. Explain the following notions and theorems (30 points): a. Decidable set/enumerable set/effectively enumerable set b. The Continuum Hypothesis/The Generalized Continuum Hypothesis c. Semantical consequence/ syntactical consequence/logical truth/theorem d. The Compactness Theorem 2. Multiple choices (each question might have more than one "correct" choice) (30 points): An open sentence ((1) may be satisfied by all sequences of objects; a. (2) may be satisfied by no sequence of objects: (3) may be satisfied by some sequences of objects yet not satisfied by others .) A closed sentence ((1) may be satisfied by all sequences of objects: (2) may be satisfied by no sequence of objects: (3) may be satisfied by some sequences of objects yet not satisfied by others.) The set of natural numbers is a proper subset of the set of rational numbers. ((1) Therefore, there isn't; (2) Still, there is) a 1-1 onto mapping from the former to the latter.) Which of the following sets of connectives is truth-functionally d. (complete: $((1) \{ \sim, \vee \}; (2) \{ \sim, \land \}; (3) \{ \sim, \supset \}; (4) \{ \sim, \equiv \}; (5) \{ \supset, \equiv \})$.) Which of the following statements is true: ((1) The union of a denumerable set and a finite set is denumerable: (2) The union of a countable set and a countable set is countable: (3) There are uncountably many truths of arithmetic.)) Which of the following statements is true: ((1) $\aleph_0 + \aleph_0 = \aleph_0$; (2) $\aleph_0 \times \aleph_0$ f. ($> \aleph_0$: (3) $2^{\aleph^0} = C$: (4) $\aleph_0 \times C > C$: (5) $\aleph_0^n > \aleph_0$. Let $I(G)=\{\langle x, y\rangle | x\in \mathbb{N} \text{ and } y\in \mathbb{N} \& x\geq y\}$, then the sequence $\langle 5, 10, 10 \rangle$ 15. 5. 10. 15. ... > satisfies ((1)(x) Gxx; (2)(x) Gxy; (3)(x)(y)Gxy; (4) $(\exists x)Gxy$).) Although first-order logic is undecidable, there is an effective positive test for ((1) satisfiability: (2) unsatisfiability: (3) validity).) If a first-order wff is satisfiable, it is satisfied in an ((1) enumerable: (2) arbitrary infinite; (3) finite) domain.) Which of the following are not decidable? ((1) sentential logic: (2) j. monadic predicate logic: (3) first-order logic: (4) arithmetic) .
- 3. Specify a separate model for each of the following items showing that (20 points):
 - a. " $(x)(y)(z)\{[(Rxy \& Ryz) \supset Rxz] \& \sim Rxx\} \& \sim (y)(\exists x)Rxy$ " is satisfiable in a model with finite domain.
 - b. It is false that $(x)(Fx \supset (\exists y)(Gy \& Rxy)) \models (\exists y)(Gy \& (x)(Fx \supset Rxy))$.

- 4. Choose one of the followings (20 points):
 - a. Prove Lindenbaum's lemma for first-order logic: If K is a consistent first-order theory, then there is a first order theory K' that is a maximal consistent extension of K with the same formulas as K.
 - b. Let ϕ be any wff, x and y any variables. <D, V> any model, and μ any value assignment to the variables. Then, where σ is just like μ except that $\sigma(x) = \mu(y)$, prove that $V_{\sigma}(\phi) = V_{\mu}(\phi[y/x])$. ($\phi[y/x]$ differs from ϕ only in that wherever ϕ has free variable x, $\phi[y/x]$ has variable y which is free in $\phi[y/x]$ but not in ϕ .
 - c. Prove: (a) $\Gamma \cup \{ \neg \phi \}$ is first-order (syntactically) inconsistent iff $\Gamma \vdash_{FOL} \phi$; (b) $\Gamma \cup \{ \phi \}$ is first-order (syntactically) inconsistent iff $\Gamma \vdash_{FOL} \neg \phi$.
 - d. Prove that the set of all subsets of natural numbers is not enumerable.