

# 國立中正大學 111 學年度碩士班招生考試試題

科目名稱：初階邏輯

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系所組別：哲學系

In this test,

“ $\neg$ ” means “not”,

“ $\wedge$ ” means “and”,

“ $\vee$ ” means “or”,

“ $\rightarrow$ ” means “if...then...”,

“ $\leftrightarrow$ ” means “if and only if”,

“ $\forall x$ ” means “for all x”, and

“ $\exists x$ ” means “for some x”.

## I. True or False

Answer each of the following questions simply by writing **True** or **False**. (24 points; 3 points each)

1.  $(A \wedge E) \rightarrow [\neg(A \leftrightarrow F) \leftrightarrow (\neg E \leftrightarrow F)]$  is a tautology.
2. If  $R(a, b)$  logically implies  $R(b, a)$ , then  $a=b$ .
3. If  $\alpha$ ,  $\beta$  and  $\gamma$  are inconsistent, then  $(\beta \leftrightarrow \gamma) \leftrightarrow \alpha$  cannot be true.
4.  $\forall x \forall y \alpha(x, y)$  is satisfiable if and only if  $\alpha(c, c)$  is satisfiable, where  $c$  is a constant symbol which does not occur in  $\alpha(x, y)$ .
5.  $\forall x \exists y \forall z (P(x, z) \rightarrow Q(y))$  is logically equivalent to  $\exists x \exists z P(x, z) \rightarrow \exists y Q(y)$ .
6. Suppose  $S$  is a set of infinitely many sentences and  $T$  is a finite subset of  $S$ . Then there must be a theorem of  $S$  which is not provable from  $T$ .
7. Suppose  $S$  and  $T$  are two consistent sets of sentences such that  $S \cup T$  is inconsistent. Then for some  $\alpha$  in  $T$ ,  $S$  proves  $\neg\alpha$ .
8. No finite model can satisfy  $\forall x \exists y (x \neq y \wedge R(x, y)) \wedge \forall x \forall y \forall z (R(x, y) \rightarrow (R(y, z) \rightarrow R(x, z)))$ .

## II. Symbolize the following two sentences. (20 points; 10 points each)

- (1) 每一對已生育的夫妻每月可領取一萬元的補助，但月收入若已達到十二萬元(含)以上則不得領取。
- (2) At least two philosophers are talking nonsense unless everyone who attends the meeting is from Mars.

**III.** Define a sentential connective  $+$  as follows.  $A+B$  is true if and only if exactly one of  $A$  and  $B$  is true (that is,  $A$  is true and  $B$  is false or  $A$  is false and  $B$  is true). Show that  $\rightarrow$  is not definable by using  $\neg$ ,  $\leftrightarrow$  and  $+$  (that is, no sentence in the language using  $A$ ,  $B$ ,  $\neg$ ,  $\leftrightarrow$  and  $+$  is equivalent to  $A \rightarrow B$ ). (10 points)

**IV.** For each of the following arguments, prove it or refute it by giving a counterexample. (46 points)

- (1) 前提 1: 並非所有的聰明人都是哲學家。 前提 2: 所有的哲學家都是苦命人。 因此, 有些聰明

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人不是苦命人。(6 points)

(2)  $\forall x(Rx \leftrightarrow Qx), \exists x(\neg(Px \leftrightarrow Qx) \leftrightarrow Rx), \forall y(S(x, y) \rightarrow Rx) \therefore \forall x((\exists yRy \wedge \exists yQy) \rightarrow Px) \rightarrow \neg S(x, b)$  (10 points)

(3)  $\neg(A \leftrightarrow B) \leftrightarrow C, B \rightarrow (C \rightarrow D) \therefore (D \wedge A) \vee (B \leftrightarrow C)$  (10 points)

(4)  $\therefore \exists z \forall y \exists x((P(z) \rightarrow R(y)) \rightarrow (P(x) \rightarrow R(x)))$  (10 points)

(5)  $\forall x \exists y R(x, y), \forall x \forall y (R(x, y) \rightarrow \exists z (z \neq x \wedge z \neq y \wedge R(x, z) \wedge R(z, y))) \therefore \exists x \exists y \exists z (x \neq y \wedge x \neq z \wedge y \neq z \wedge R(x, y) \wedge R(y, z))$  (10 points)